

Last time: scissors congruence.

Recall: If $P, Q \in \mathbb{R}^3$ are two polyhedra such that

$D_f(P) \neq D_f(Q)$ for some $\mathbb{Q}$ -linear function f with $f(\mathbf{n})=0$
then P and Q are not scissors congruent.

What about the converse?

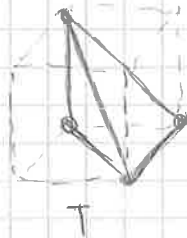
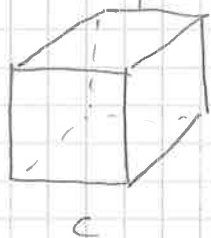
Theorem (Sydler)

Polyhedra P and Q are scissors congruent if
 $D_f(P) = D_f(Q)$ for every $\mathbb{Q}$ -linear function f with $f(\mathbf{n})=0$.

Proof: is quite involved.

Example / Exercise

- Show that for the unit cube C and the tetrahedron T below



$$D_f(C) = D_f(T) = 0$$

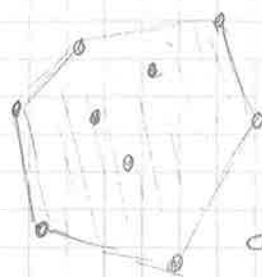
for every f with $f(\mathbf{n})=0$.

As a consequence, T is scissors congruent to a cube (of volume $1/6$)

- Try to find a dissection of T and reassemble the pieces to form a cube.

Today: Convex hulls

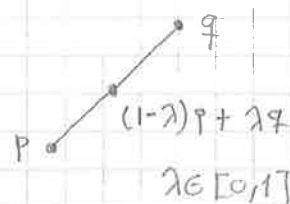
Before defining the convex hull, we start with some basic definitions.



For $S \subseteq \mathbb{R}^d$
 $\text{conv}(S)$

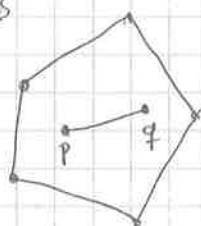
For $p, q \in \mathbb{R}^d$, the line segment between p and q is:

$$\overline{pq} := \{(1-\lambda)p + \lambda q : \lambda \in [0, 1]\}$$

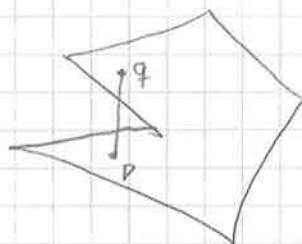


Def. A set $S \subseteq \mathbb{R}^d$ is convex if for every $p, q \in S$, the line segment $\overline{pq} \subseteq S$.

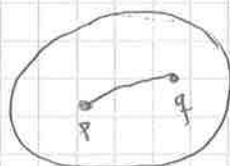
Examples



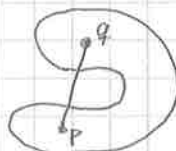
convex



not convex



convex



not convex

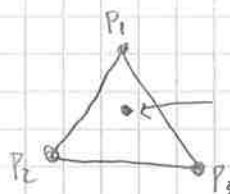
Lemma If $\{S_i\}_{i \in I}$ are convex, then

$\bigcap_{i \in I} S_i$ is also convex

Proof For $p, q \in \bigcap S_i$, the segment $\overline{pq} \subseteq S_i \Rightarrow \overline{pq} \subseteq \bigcap S_i$

For $p_1, p_2, p_3 \in \mathbb{R}^d$, the triangle with vertices p_1, p_2, p_3 is

$$\{\lambda_1 p_1 + \lambda_2 p_2 + \lambda_3 p_3 : \lambda_i \geq 0, \lambda_1 + \lambda_2 + \lambda_3 = 1\}$$



$$\lambda_1 p_1 + \lambda_2 p_2 + \lambda_3 p_3$$

$$\lambda_i \geq 0, \lambda_1 + \lambda_2 + \lambda_3 = 1$$

In general:

Def (Convex hull)

For $S \subseteq \mathbb{R}^d$, the convex hull of S is

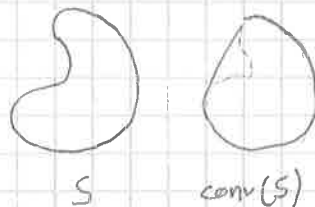
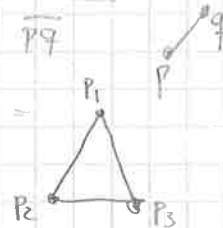
$$\text{conv}(S) := \left\{ \sum_{i=1}^n \lambda_i s_i : n \in \mathbb{N}, s_i \in S, \lambda_i \geq 0, \sum_{i=1}^n \lambda_i = 1 \right\}$$

↳ convex linear combination

Examples:

• $S = \{p, q\} \Rightarrow \text{conv}(S) = \overline{pq}$

• $S = \{p_1, p_2, p_3\} \Rightarrow \text{conv}(S) =$



Remarks: The following hold:

- $S \subseteq \text{conv}(S)$ ✓
- $S' \subseteq S \Rightarrow \text{conv}(S') \subseteq \text{conv}(S)$ ✓

Lemma: The convex hull $\text{conv}(S)$ is convex.

Proof Let $p, q \in \text{conv}(S)$.
Want to show that $\overline{pq} \subseteq \text{conv}(S)$

Without loss of generality

$$p = \sum_{i=1}^n \lambda_i s_i, \quad q = \sum_{i=1}^n \mu_i s_i \quad \lambda_i, \mu_i \geq 0 \quad \sum_{i=1}^n \lambda_i = \sum_{i=1}^n \mu_i = 1$$

for some set $\{s_1, \dots, s_n\} \subseteq S$

For $\alpha \in [0, 1]$

$$(1-\alpha)p + \alpha q = \sum_{i=1}^n \underbrace{(\lambda_i(1-\alpha) + \mu_i\alpha)}_{w_i} s_i$$

$$\begin{aligned} \text{and } \sum_{i=1}^n (\lambda_i(1-\alpha) + \mu_i\alpha) &= (1-\alpha) \sum_{i=1}^n \lambda_i + \alpha \sum_{i=1}^n \mu_i \\ &= (1-\alpha) + \alpha \\ &= 1 \end{aligned}$$

Lemma $S \subseteq \mathbb{R}^d$ is convex iff $\text{conv}(S) = S$

Proof: ($\Leftarrow$) Assume $\text{conv}(S) = S$.

Let $p, q \in S$, then $(1-\lambda)p + \lambda q \in \text{conv}(S) = S$

So, S is convex.

($\Rightarrow$) Assume S is convex. We already know that $S \subseteq \text{conv}(S)$.

Let $p = \sum_{i=1}^n \lambda_i s_i \in \text{conv}(S)$ for some $s_1, s_2, \dots, s_n \in S$
 $\lambda_i \geq 0$
 $\sum \lambda_i = 1$

We want to show that $p \in S$.

We proceed by induction on n .

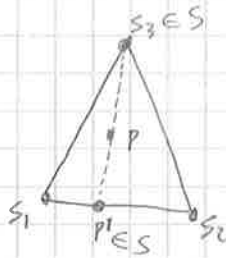
$n=1$ ✓

Assume $n \geq 2$.

Let $\lambda' = \sum_{i=1}^{n-1} \lambda_i$

and

$p' = \sum_{i=1}^{n-1} \frac{\lambda_i}{\lambda'} s_i \in S$ by induction hypothesis.



Idea: • $p' \in S$ by induction
 • $p \in S$ by convexity

Moreover, since $\lambda_n = 1 - \lambda'$ we have

$$p = \underbrace{\lambda' p'}_{\in S} + (1 - \lambda') \underbrace{s_n}_{\in S}$$

Therefore $p \in S$ by convexity of S . $\blacksquare$

Theorem $\text{conv}(S) = \bigcap_{\substack{C \supset S \\ C \text{ convex}}} C$

Proof

($\supseteq$) The set $\text{conv}(S)$ is convex and contains S .

$$\text{Thus } \bigcap_{\substack{C \supset S \\ C \text{ convex}}} C \subseteq \text{conv}(S)$$

($\subseteq$) If $S \subseteq C$ and C is convex, then

$$\text{conv}(S) \subseteq \text{conv}(C) = C.$$

Therefore

$$\text{conv}(S) \subseteq \bigcap_{\substack{C \supset S \\ C \text{ convex}}} C$$

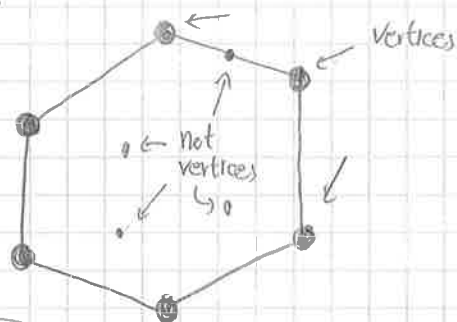
Corollary $\text{conv}(S)$ is the smallest convex set containing S .

Vertices of convex hulls

Let $P \subseteq \mathbb{R}^d$ be a finite collection of points.

$p \in P$ is called a vertex of $\text{conv}(P)$

If $p \notin \text{conv}(P_1, P_3)$



Exercise

Let $V \subseteq P$ be the subset of vertices of P .

Then

$$\text{conv}(V) = \text{conv}(P)$$