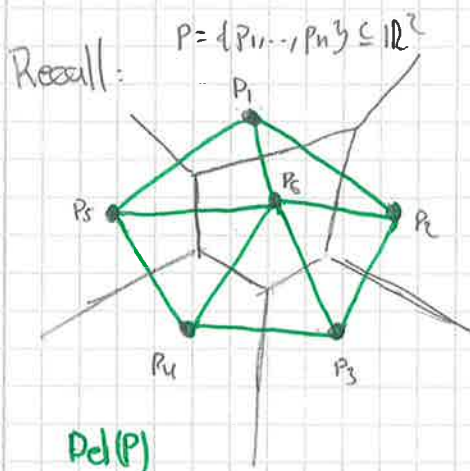


Last time :

Delaunay graph

Today :

Delaunay triangulations
The parabolic lift



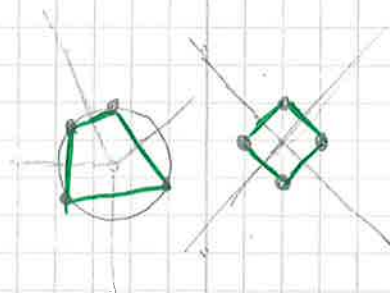
The Delaunay graph

$\text{Del}(P)$ is the straight-line embedding of the dual graph of the Voronoi diagram $\text{Vor}(P)$

We have the following duality dictionary:

Duality:

$\text{Del}(P)$	$\text{Vor}(P)$
Vertices $p_1, \dots, p_n$	Regions $\mathcal{V}(p_1), \dots, \mathcal{V}(p_n)$
Edges $\overline{p_i p_j}$	Edges defined by bisector L_{ij}
Faces k-gon $\text{conv}\{p_{i_1}, \dots, p_{i_k}\}$	Vertices point q shared by regions $\mathcal{V}(p_{i_1}), \dots, \mathcal{V}(p_{i_k})$



Rephrasing the characterization theorem for vertices and edges of $\text{Vor}(P)$:

Theorem Let $P = \{p_1, \dots, p_n\}$ be a set of points in the plane. Then

- (i) Three points $p_i, p_j, p_k \in P$ are vertices of the same face of $\text{Del}(P)$ iff the circle through p_i, p_j, p_k has no point of P in its interior.
- (ii) Two points $p_i, p_j \in P$ form an edge in $\text{Del}(P)$ iff there is a closed disk that contains p_i, p_j on its boundary and does not contain any other point of P .

- No four points on a circle

Observe that if no four points of P lie on a circle (general position) then all the vertices of $\text{Vor}(P)$ have degree 3.

So, all the faces of $\text{Del}(P)$ are triangles.

In this case

$\text{Del}(P)$ is called the Delaunay triangulation of P .

- General case

In the general case, we have to be a bit more careful.

We define a Delaunay triangulation to be any triangulation obtained by adding edges to the Delaunay graph $\text{Del}(P)$.

(since all faces of $\text{Del}(P)$ are convex polygons, this is easy)

↓
with vertices on a circle.

Theorem Let $P = \{p_1, \dots, p_n\}$ be a set of points in the plane.

A triangulation T of P is a Delaunay triangulation iff the circumcircle of any triangle does not contain any point of P in its interior.

Theorem A triangulation T of P is legal iff T is a Delaunay triangulation of P .

Proof (Exercise).

Corollary If no four points of P lie on a circle (general position)

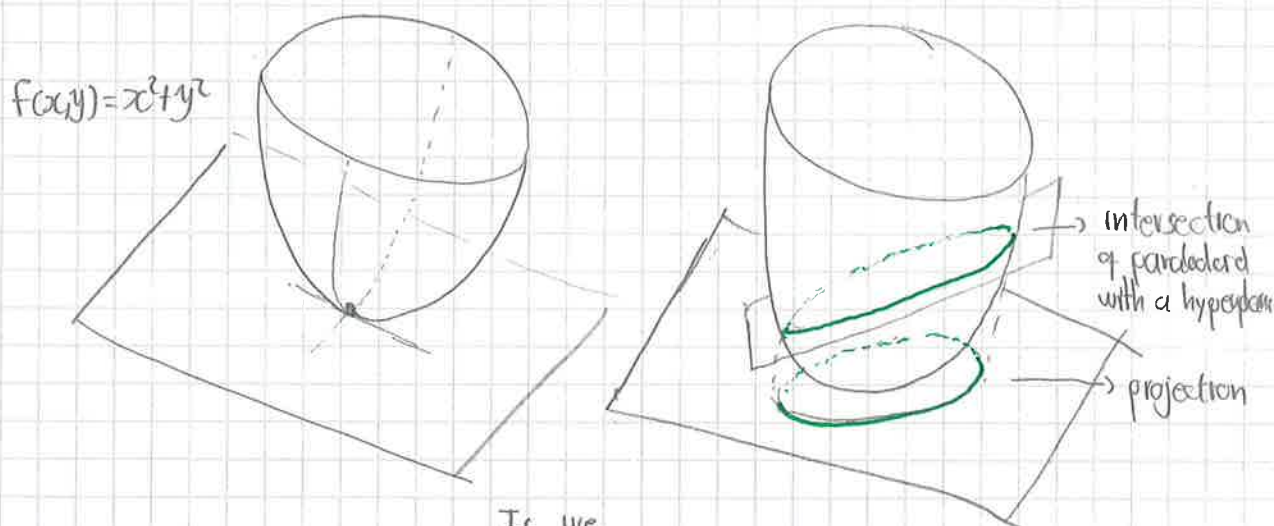
Then the Delaunay triangulation $\text{Del}(P)$ is the unique legal triangulation of P . Hence, it is the unique angle-optimal triangulation.

The parabolic lift

Goal: Obtain the Delaunay triangulation (or subdivision determined by the Delaunay graph for points not in general position) as the projection as the lower hull of a polytope in one higher dimension.

Let $P = \{p_1, \dots, p_n\}$ be a set of points in the plane

Consider the paraboloid $z = x^2 + y^2$



Interesting property of the paraboloid:

If we intersect it with a plane and project the resulting curve to the xy -plane then we obtain a circle!

Moreover, the points in the paraboloid above the plane project outside the circle, and the points below the hyperplane project inside the circle. (see proof of theorem below)

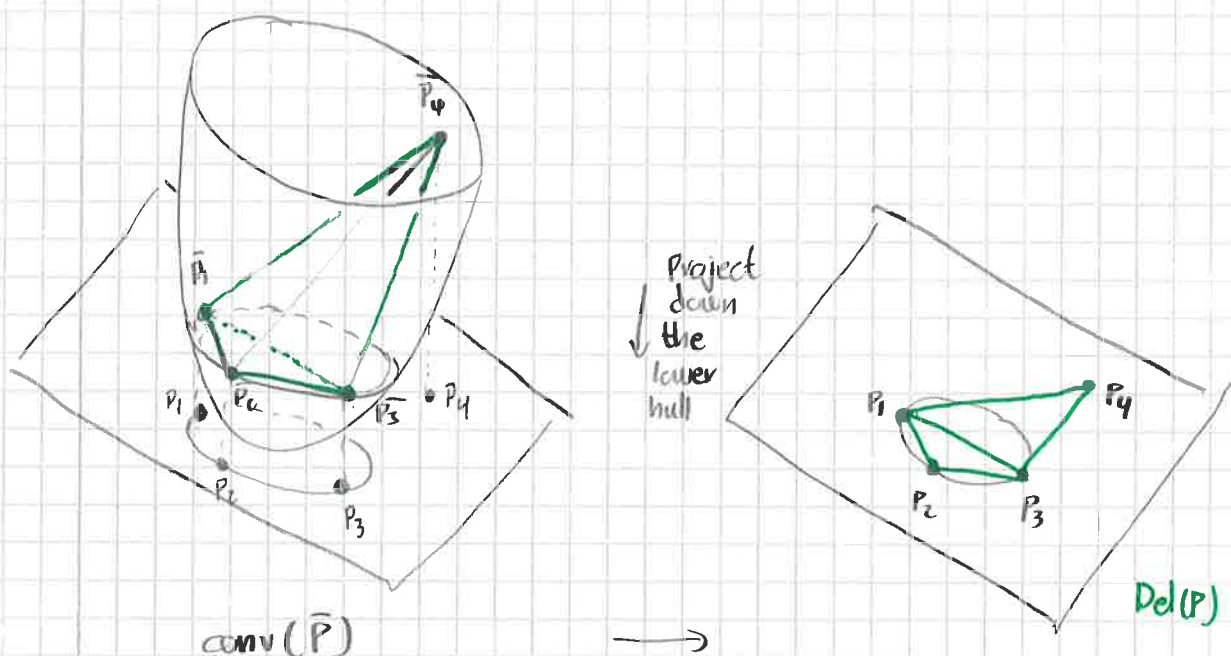
For $p = (x, y) \in P$, let $\bar{p} = (x, y, x^2 + y^2)$ be the ^{corresponding} lifted point.

Let $\bar{P} = \{\bar{p}_1, \bar{p}_2, \dots, \bar{p}_n\}$ and $\text{conv}(\bar{P})$ be the convex hull of the lifted points in $\mathbb{R}^3$ (a 3-dimensional polytope)

The lower hull of $\text{conv}(\bar{P})$ is the set of faces of $\text{conv}(\bar{P})$ that "you can see from below", i.e. such that $\text{conv}(\bar{P})$ is above the face

Theorem

The subdivision determined by the Delaunay graph $\text{Del}(P)$ (The Delaunay triangulation if no four points are in circle) is the projection of the lower hull of $\text{conv}(\bar{P})$



Proof

Idea: projection of the intersection of the paraboloid with a plane is a circle in the xy -plane

In order to show this, consider the tangent plane of the paraboloid at the point $(x_0, y_0, x_0^2 + y_0^2)$

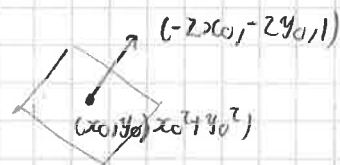
Two vectors spanning this plane are

$$\left\langle \frac{\partial}{\partial x_0} \right\rangle \rightarrow (1, 0, 2x_0)$$

$$\left\langle \frac{\partial}{\partial y_0} \right\rangle \rightarrow (0, 1, 2y_0)$$

The normal vector is

$$(-2x_0, -2y_0, 1)$$



The equation of the tangent plane is

$$-2x_0(x - x_0) - 2y_0(y - y_0) + (z - x_0^2 - y_0^2) = 0$$



$$z = 2x_0x + 2y_0y - x_0^2 - y_0^2$$

Now, a plane E intersecting the paraboloid is parallel to some tangent plane.

Thus the equation of E is of the form

$$z = 2x_0x + 2y_0y - x_0^2 - y_0^2 + r^2$$

The intersection of E with the paraboloid $z = x^2 + y^2$ is

$$x^2 + y^2 = 2x_0x + 2y_0y - x_0^2 - y_0^2 + r^2$$

$\Updownarrow$

$$\boxed{(x - x_0)^2 + (y - y_0)^2 = r^2}$$

the equation of
a circle

Hence, the projection of the faces of the lower hull of $\text{conv}(\bar{P})$ satisfy the property that

p_i, p_j, p_k are a ^(projected) face

$\Leftrightarrow$

the circle through p_i, p_j, p_k
has no point of P in its interior

Thus the projection of the lower hull is the subdivision determined by the Delaunay graph $\text{Del}(P)$.

In the case of general position (no four points of P in a circle) this is the ^{unique} Delaunay triangulation.

■

Algorithmic Consequence

Since convex hulls in 3D can be computed in $O(n \log n)$ time (for example using a Divide-and-Conquer algorithm - Exercise).

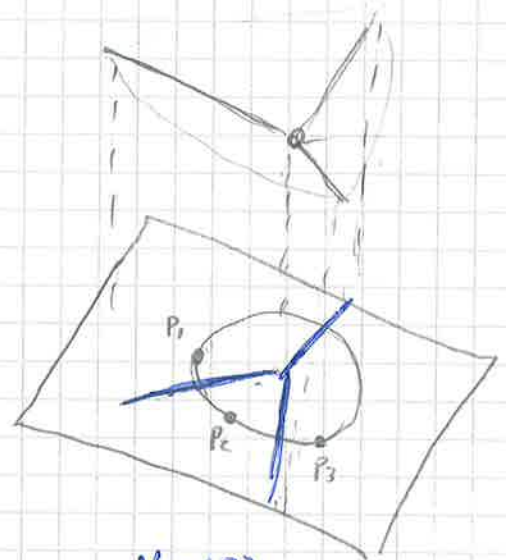
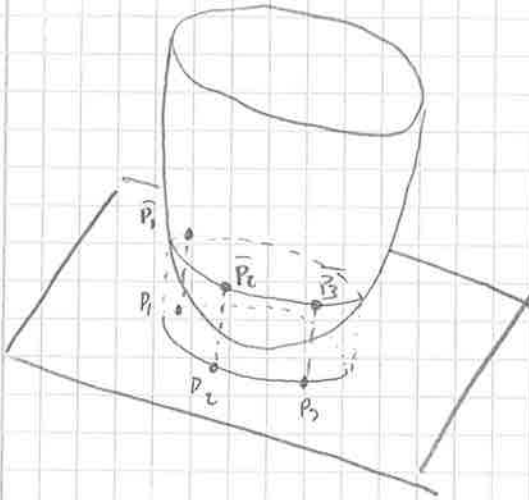
Then, we can compute Delaunay triangulations in $O(n \log n)$ time (as well as Voronoi diagrams).

Remark: Duality between Delaunay triangulations and convex hulls holds in higher dimensions.

Application: Solid meshing of 3D objects using Delaunay tetrahedralizations.

Questions

How can we obtain the Voronoi diagram using the parabolic lift?



$\text{Vor}(P)$

Answer

Consider the tangent planes of the paraboloid at the lifted points $\bar{P}_1, \dots, \bar{P}_n$.

Look at these hyperplanes and their intersections from above. ($z = +\infty$)

The projection of the upper envelope (pieces not obscured by any other plane) yield the Voronoi diagram.